

## Relations among Embedding Parameters for Graphs

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### ABSTRACT

Let  $G$  be a simple graph with  $n$  vertices, and let  $a_1$ ,  $bt$ ,  $\gamma$ , and  $\theta$  denote, respectively, the arboricity, book thickness (also called pagewidth), genus, and thickness of  $G$ . We establish inequalities, each of which is best possible up to a constant, between pairs of these parameters. In particular, we show:

$$\begin{aligned}\theta &\leq 6 + \sqrt{2\gamma - 2}, \\ a_1 &\leq bt + 1, \quad \text{and} \\ \gamma &\leq (\theta - 1)(n - 1).\end{aligned}$$

### 1. Introduction

There are several different ways to characterize the embeddability of a graph  $G$ . Some of the numerical measures include the genus  $\gamma$ , the thickness  $\theta$ , the book thickness (or pagewidth)  $bt$ , and the arboricity  $a_1$ . These parameters have many applications, including VLSI design and complexity theory.

In this paper, we study relationships between these parameters. We prove several inequalities, each of which is best possible up to a constant.

### 2. Definitions and Basic Properties

Throughout,  $G$  will denote a simple, *connected* graph, and  $n$  and  $e$  (or  $n_G$  and  $e_G$ , if more than one graph is being considered) will denote, respectively, the number of

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vertices and edges of  $G$ . We will consider four parameters that measure the embeddability of  $G$  in different ways.

First, the (orientable) *genus of  $G$* , denoted by  $\gamma$ , is the minimum number  $k$  such that  $G$  can be embedded on a sphere with  $k$  handles. Second, the *thickness of  $G$* , denoted by  $\theta$ , is the minimum number of edge-disjoint planar subgraphs whose union is  $G$ . Third, the *arboricity of  $G$* , denoted by  $a_1$ , is the minimum number of edge-disjoint forests whose union is  $G$ . Finally, the *book thickness or pagewidth of  $G$* , denoted by  $bt$ , is the minimum number of pages joined at a common spine such that  $G$  can be embedded with its vertices on the spine and edges on the pages.

All these parameters are known for the complete graphs, the complete bipartite graphs, and the generalized cubes. For example,  $\gamma(K_n) = \lceil (n-3)(n-4)/12 \rceil$  for  $n \geq 3$ ,  $\theta(K_n) = \lfloor (n+7)/6 \rfloor$  (except that  $\theta(K_9) = \theta(K_{10}) = 3$ ),  $a_1(K_n) = \lceil n/2 \rceil$ , and  $bt(K_n) = \lceil n/2 \rceil$  for  $n \geq 4$  (see [BCL, BK]).

There are two fundamental results concerning these parameters that will be used throughout the paper. The first is Euler's formula (and a corollary), and the second is Nash-Williams' formula for arboricity:

**Theorem 1** *If  $G$  is embedded on  $S_g$ , the sphere with  $g$  handles, then, letting  $f$  denote the number of faces:*

- (i)  $n - e + f \geq 2 - 2g$ , with equality if and only if the embedding is a 2-cell embedding,
- (ii)  $e \leq 3n - 6 + 6g$ .

**Theorem 2 [N]** *For any graph  $G$ ,*

$$a_1 = \max_{H < G} \left\lceil \frac{e_H}{n_H - 1} \right\rceil,$$

where  $H$  ranges over all nontrivial induced subgraphs of  $G$ .

It follows easily from Theorem 2 that  $a_1$  is bounded in terms of the minimum degree  $\delta$  and maximum degree  $\Delta$  of  $G$ :

**Corollary 2** (i)  $\left\lceil \frac{\delta+1}{2} \right\rceil \leq a_1 \leq \left\lceil \frac{\Delta+1}{2} \right\rceil$ .

(ii) [AH] *If  $G$  is  $d$ -regular, then  $a_1 = \left\lceil \frac{d+1}{2} \right\rceil = \left\lceil \frac{e}{n-1} \right\rceil$ .*

The arboricity and thickness of a graph  $G$  are closely related in size; namely,  $\theta \leq a_1 \leq 3\theta$ . The first inequality follows from the definitions, and the second follows from Theorems 1 and 2, which together imply that any planar graph has arboricity at most 3. Furthermore, these bounds are sharp, in the sense that for any  $\theta$ , there are graphs that achieve each extreme; see [T, BH].

Another related parameter is the *vertex arboricity*, denoted by  $a$ . The vertex arboricity of a graph  $G$  is the minimum number of vertex-disjoint forests whose union is  $G$ . In [B] it is shown that  $a \leq a_1$ , and so any bound on arboricity will give the same bound for vertex arboricity.

### 3. Upper Bounds on Thickness and Arboricity

In [A], Asano showed that triangle-free graphs and toroidal graphs have thickness  $\theta \leq \gamma + 1$ . A recent result of Scheinerman [S] shows that for any  $G$ , the arboricity  $a_1$  satisfies  $a_1 \leq 2 + \sqrt{3\gamma}$ . Since  $\theta \leq a_1$ , it follows that, for any  $G$ ,  $\theta \leq 2 + \sqrt{3\gamma}$ . Archdeacon and Richter (unpublished manuscript) have obtained a similar bound on  $\theta$ . Furthermore, the example of the complete graph  $K_n$  shows that this result is best possible up to a constant, since  $\theta(K_n) \approx \sqrt{\gamma(K_n)} / 3$ .

Using an idea due to M. Albertson and C. Thomassen (personal communication), we improve on Scheinerman's result for thickness, reducing the constant from  $\sqrt{3}$  to  $\sqrt{2}$ .

**Theorem 3** *If  $G$  is a simple graph, then*

$$\theta \leq 6 + \sqrt{2\gamma - 2}.$$

**Proof** Let  $k = 6 + \lfloor \sqrt{2\gamma - 2} \rfloor$  and let  $G_1, G_2, \dots, G_k$  each be graphs with the same vertex set as  $G$ , but initially no edges. For each vertex  $v$  with  $\deg(v) \leq k$ , remove its edges from  $G$  and place one edge in each  $G_i$  for  $1 \leq i \leq \deg(v)$ . Repeat this process for each vertex of degree at most  $k$  in the resulting graph until all the remaining vertices in the final graph  $G'$  have degree strictly greater than  $k$ .

Each of the graphs  $G_i$  generated by this process is acyclic: Suppose some  $G_i$  contained a cycle  $C$ . Consider the "oldest" vertex  $v$  on the cycle  $C$ , i.e., the vertex on  $C$  whose edges were removed from  $G$  before any other vertices of  $C$ . Then by the construction of the  $G_i$ 's, the degree of  $v$  in  $C$  is 0 or 1, a contradiction.

Since all remaining vertices of  $G'$  have degree at least  $k + 1 \geq 6 + \sqrt{2\gamma - 2}$ , Thm. 1 (ii) yields

$$(6 + \sqrt{2\gamma(G) - 2}) n' \leq 2e' \leq 6n' + 12(\gamma(G') - 1),$$

and, since  $\gamma(G') \leq \gamma(G)$ ,

$$(6 + \sqrt{2\gamma(G) - 2}) n' \leq 2e' \leq 6n' + 12(\gamma(G) - 1),$$

where  $n'$  and  $e'$  are the number of vertices and edges of  $G'$ . Thus,

$$n' \leq 6\sqrt{2\gamma(G) - 2}.$$

Now  $G'$  is a subgraph of the complete graph  $K_{n'}$ , which satisfies

$$\theta(K_{n'}) = \begin{cases} \lfloor (n' + 7) \rfloor / 6, & n' \neq 9, 10 \\ 3, & n' = 9, 10 \end{cases}$$

It follows that

$$\begin{aligned}\theta(G') &\leq \theta(K_{n'}) \\ &\leq \lfloor (n' + 9) / 6 \rfloor \\ &\leq \lfloor (6\sqrt{2\gamma(G)} - 2 + 9) / 6 \rfloor \\ &\leq \sqrt{2\gamma(G)} - 2 + 6.\end{aligned}$$

Hence,  $G'$  can be partitioned into at most  $k = \lfloor \sqrt{2\gamma(G)} - 2 \rfloor + 6$  planar subgraphs, say  $P_1, \dots, P_{k'}$ , where  $1 \leq k' \leq k$ . Form the unions  $P_i \cup G_i$  for  $i = 1, \dots, k'$ .

We claim that each tree of  $G_i$  is incident with at most one vertex of  $P_i$  and hence  $P_i \cup G_i$  is a planar graph. If not, find a tree  $T_i$  of  $G_i$  and two vertices  $x$  and  $y$  of  $P_i \cap T_i$  such that the path  $P_{xy}$  from  $x$  to  $y$  in  $T_i$  contains no other vertex of  $P_i$ . Note that  $x$  and  $y$  are not adjacent since each edge of  $G_i$  is incident with at least one vertex not in  $P_i$ . As before, let  $z$  be the "oldest" internal vertex of  $P_{xy}$ . Its degree in  $P_{xy}$  is 0 or 1, a contradiction.

In summary, the original graph  $G$  has thickness at most  $k$ , where  $k \leq 6 + \sqrt{2\gamma(G)} - 2$ .  $\square$

Since  $a_1 \leq 3\theta$ , we obtain as an immediate corollary that  $a_1(G) \leq 3\sqrt{2\gamma(G)} - 2 + 18$ . However, we can improve the constant if, in the proof of Theorem 3, we replace "thickness" by "arboricity" and use the arboricity of  $K_{n'}$  to estimate  $a_1(G')$ . The resulting inequality is weaker than Scheinerman's best possible result that  $a_1(G) \leq 2 + \sqrt{3\gamma(G)}$  (cf. [S]):

**Corollary 3** *If  $G$  is a simple graph, then*

$$a_1(G) \leq 6 + \sqrt{6\gamma(G)} - 6.$$

The examples of the complete graph, the complete bipartite graph, and the generalized cube all have  $\theta = 2\gamma/n$ , and this suggests a natural bound on  $\theta$  in terms of  $\gamma$  and  $n$ . Since for all graphs  $\gamma \leq n^2/12$ , we have  $\gamma/n \leq \sqrt{\gamma}$ , suggesting a tighter upper bound:  $\theta = O(\gamma/n)$ . However, this bound does not hold in general: Let  $G$  be any graph with  $\theta = O(\gamma/n)$ , and embed  $G$  on its genus surface. Append an arbitrary number of vertices of degree one to  $G$  and note that the resulting graphs may no longer satisfy a  $O(\gamma/n)$  bound, since  $n$  can become arbitrarily large. In contrast, in the next section we will show that  $\gamma = O(\theta n)$  for all graphs.

However, for many graphs including the regular graphs it is true that  $a_1 = O(\gamma/n)$  and hence  $\theta = O(\gamma/n)$ .

**Theorem 4** *Let  $0 < \varepsilon \leq 1$ . If  $G$  is a graph with at least  $1 + 1/\varepsilon$  vertices and such that  $a_1(G) = \lceil e / (n - 1) \rceil$ , then*

$$a_1(G) \leq \lceil (3 + 6\gamma/n)(1 + \varepsilon) \rceil.$$

**Proof** By Theorem 1 (ii),

$$\begin{aligned} a_1(G) &= \lceil e / (n - 1) \rceil \\ &\leq \lceil (3n - 6 + 6\gamma) / (n - 1) \rceil \\ &\leq \lceil (3 + 6\gamma/n) (1 + 1 / (n - 1)) \rceil \\ &\leq \lceil (3 + 6\gamma/n) (1 + \epsilon) \rceil. \quad \square \end{aligned}$$

From Corollary 2 (ii) we have:

**Corollary 4** *Let  $0 < \epsilon \leq 1$ . If  $G$  is a regular graph with at least  $1 + 1/\epsilon$  vertices, then*

$$a_1(G) \leq \lceil (3 + 6\gamma/n) (1 + \epsilon) \rceil = O(\gamma/n).$$

Since  $\theta \leq a_1$ , we get the same bound on  $\theta$ .

In [H], Hutchinson shows that a graph  $G$  with a 2-cell embedding on the  $k$ -handled sphere  $S_k$  is 5-colorable if the embedding has "short" edges. There, the length of an edge is measured relative to the length of a side of the standard  $4k$ -gon on which the embedding is represented. The following theorem shows that under those conditions,  $G$  has thickness at most 2 (and hence arboricity at most 6).

**Theorem 5** *Suppose  $G$  has a 2-cell embedding on a surface of genus  $k \geq 1$ , and suppose that  $G$  has a representation  $G_k$  on the standard  $4k$ -gon (i.e., a  $4k$ -gon in which each side has unit length) such that every edge of  $G_k$  has length  $< \epsilon$ , where  $\epsilon$  is at most  $1/2$ . Then  $\theta(G) \leq 2$ .*

**Proof** We begin by extending  $G_k$  to a triangulation of the surface. Subdivide every nontriangular face by adding a new vertex adjacent to all vertices on the face boundary. If any new edge has length  $\epsilon$  or more, subdivide it by adding vertices. Repeat this process until the resulting graph  $G_k'$  is a triangulation of the surface having all edges of length less than  $\epsilon$ .  $G_k$  is a subgraph of  $G_k'$  which implies that  $\theta(G) \leq \theta(G')$ . Thus, we may assume that  $G_k$  was a triangulation to begin with.

For each  $i = 1, \dots, k$ , consider the "handle" in the polygon  $P_k$  with sides labelled  $a_i, b_i, a_i^{-1}, b_i^{-1}$ . Note that one point  $S$  is common to all  $4k$  sides of  $P_k$ . Let  $p_i$  (resp.,  $q_i$ ) be the set of all points of  $P_k$  at distance  $2\epsilon$  (resp.,  $\epsilon$ ) from side  $b_i$ . Note that  $\epsilon \leq 1/2$  and so  $2\epsilon \leq 1$ . Thus,  $p_i$  and  $q_i$  are paths going from  $a_i$  to  $a_i^{-1}$  representing a non-null-homologous cycle (henceforth abbreviated "nnh-cycle") on the surface (but not necessarily in the graph). Let  $L_i$  be the closed region bounded by  $p_i, q_i, a_i$ , and  $a_i^{-1}$ , i.e.,  $L_i$  is the set of all points of  $P_k$  to the left of  $p_i$  (as we travel from  $a_i$  to  $a_i^{-1}$ ) and within distance  $\epsilon$  of  $p_i$  (see Fig. 1 below).

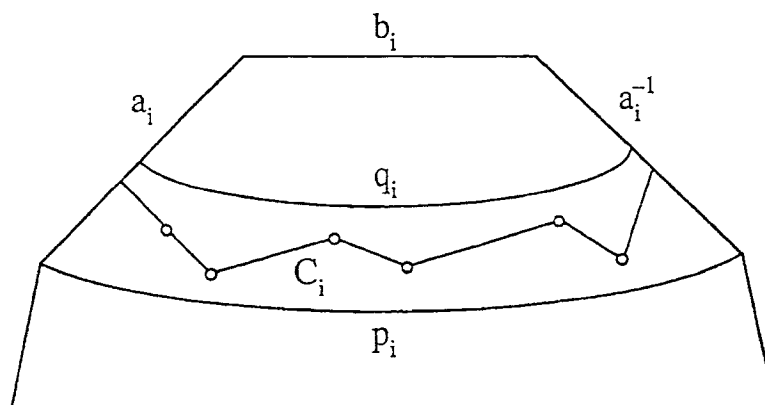


Figure 1

Within  $L_i$ , there is a path, starting at corresponding edges or vertices on  $a_i$  and  $a_i^{-1}$ , that represents a  $nnh$ -cycle in  $G_k$ : To find such a path, color a face or region of  $G_k$  blue if it meets  $L_i$ , but does not cross  $p_i$ . Color the remaining regions red. Because  $G_k$  is a 2-cell triangulation and the width of  $L_i$  is  $\varepsilon$ , we can only cross red faces as we walk along the path  $q_i$ . Similarly, we can only cross blue faces as we walk along  $p_i$ . Therefore, the boundary between the red and blue regions must lie within  $L_i$  and it must reach from  $a_i$  to  $a_i^{-1}$ , giving us a  $nnh$ -cycle in  $G$ .

For each  $i = 1, \dots, k$ , choose such a cycle  $C_i$ , and let  $H_i$  be the subgraph of  $G$  induced by the set of all edges incident to a vertex of  $C_i$  and lying to the left of  $C_i$ . Since all vertices of  $C_i$  lie within  $L_i$ , no vertex of  $H_i$  can lie on  $b_i$ . Thus  $H_i$  is planar, since it is embedded on a cylinder. Since, by construction, the  $H_i$ 's are vertex-disjoint, their union  $H$  is also planar.

Now let  $K$  be the subgraph of  $G$  induced by the edges of  $G - H$ . Then  $K$  is planar because it is embedded on the surface that remains when we cut  $P_k$  along all the  $nnh$ -cycles  $C_i$ , i.e.,  $K$  is embedded on a sphere with open disks cut out and is therefore planar. Finally,  $G = H \cup K$ , and so  $\theta(G) \leq 2$ , as claimed.  $\square$

It follows easily from the definitions that  $\theta \leq \lceil bt/2 \rceil$ , where  $bt$  is the book thickness of  $G$  (also cf. [BK]). This in turn implies that  $a_1 \leq 3\lceil bt/2 \rceil$ . In fact,  $a_1 \leq bt + 1$ , as shown in the next theorem.

**Theorem 6** *If  $G$  is a simple graph, then*

$$a_1 \leq bt + 1.$$

**Proof** The inequality clearly holds if  $a_1 = 1$ , so assume that  $a_1 \geq 2$ . By [BK Thm. 3.3], the following inequality holds:

$$bt \geq (e - n) / (n - 3), \text{ or equivalently,}$$

$$e \leq n(bt + 1) - 3bt.$$

Let  $H$  be a subgraph of  $G$  that gives the maximum value in the Nash–Williams arboricity formula. Then

$$a_1 = \lceil e_H / (n_H - 1) \rceil$$

$$\leq \lceil [n_H(bt_H + 1) - 3bt_H] / (n_H - 1) \rceil$$

$$\leq \lceil bt_H + 1 + (1 - 2bt_H) / (n_H - 1) \rceil.$$

Since  $2 \leq a_1 = \lceil e_H / (n_H - 1) \rceil$ ,  $H$  is not a path and hence  $bt_H \geq 1$ . Thus,  $1 - 2bt_H$  is negative, and therefore  $a_1 \leq bt_H + 1 \leq bt + 1$ .  $\square$

**Example 1** A variation on a construction in [BK] provides an infinite class of graphs whose arboricity achieves this upper bound: For  $k \geq 3$ , fix  $n \geq 2k + 1$  and let  $T_k$  be the triangulated  $n$ -gon shown in Figure 2.

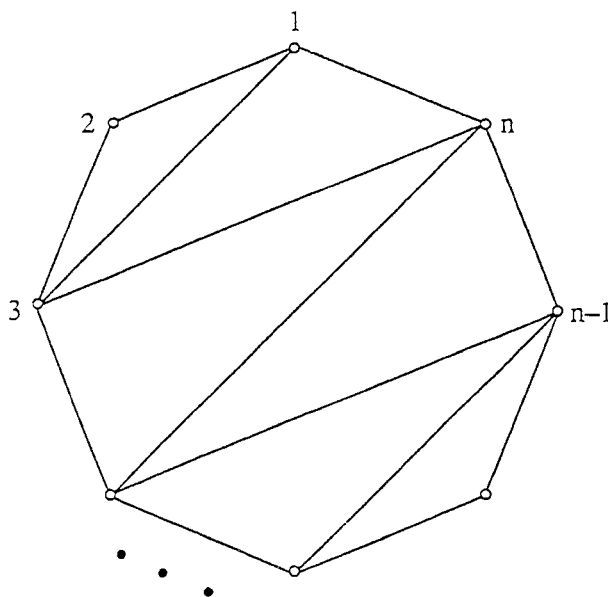


Figure 2

For  $i = 0, \dots, k - 1$ , let  $G_k^i$  be the graph obtained by rotating  $T_k$  through  $i$  successive positions (but leaving vertex labels fixed). Let  $G_k$  be the graph of order  $n$  obtained by taking the union of the  $G_k^i$  for  $i = 0, \dots, k - 1$ . Clearly,  $\delta(G_k) \leq k$ , and so  $a_1(G_k) \leq k + 1$  by Theorem 5. Thus, it suffices to show that  $a_1(G_k) \geq k + 1$ .

By construction of  $G_k$ , we have

$$\begin{aligned} e &= k(n - 3) + n \\ &= (k + 1)n - 3k. \end{aligned}$$

Next,

$$\begin{aligned} a_1 &\geq \lceil e / (n - 1) \rceil \\ &= \lceil [(k + 1)n - 3k] / (n - 1) \rceil \\ &= \lceil [k + 1 - (2k - 1) / (n - 1)] \rceil. \end{aligned}$$

Since  $n \geq 2k + 1$ , it follows that  $a_1 \geq k + 1$ , and thus  $a_1 = \delta + 1$ , as claimed.

#### 4. Upper Bounds on Genus

It is impossible to find an upper bound for the genus  $\gamma$  in terms of thickness  $\theta$  or arboricity  $a_1$ , as is demonstrated by the example of  $K_{4,n}$ , which has  $\theta = 2$  and  $a_1 = \delta = 4$  for  $n \geq 4$ , but which has unbounded genus  $\gamma = \lceil (n - 2) / 2 \rceil$  (see [BCL]). The genus of a graph  $G$  can be bounded, however, in terms of both  $\theta$  and the number of vertices  $n$  (or in terms of  $a_1$  and  $n$ , or in terms of  $\delta$  and  $n$ ), as shown in the following theorems. As in Theorem 3 and its corollary, the complete graph  $K_n$  shows that these results are best possible up to a constant.

**Theorem 7** *If  $G$  is a simple graph with  $n$  vertices, then*

$$\gamma(G) \leq (\theta(G) - 1)(n - 1).$$

**Proof** Suppose  $G$  can be written as the union of  $G_1, \dots, G_\theta$ , where the  $G_i$ 's are planar and pairwise edge-disjoint. Embed each  $G_i$  on a sphere  $S_i$ ; we will sew the spheres together in such a way that the resulting surface has  $(\theta - 1)(n - 1)$  handles.

On sphere  $S_1$ , cut out  $\theta - 1$  tangent disks at each vertex  $i$ . Call the disks  $D_i^{1,2}, D_i^{1,3}, \dots, D_i^{1,\theta}$ . On each other sphere  $S_k$ , for  $k = 2, \dots, \theta$ , cut out only one tangent disk at each vertex  $i$  and call the disk  $D_i^k$ .

Join the boundaries of the cut-out disks with cylinders as follows: For  $k = 2, \dots, \theta$ , attach a cylinder from  $D_i^{1,k}$  (the  $k^{\text{th}}$  tangent disk to vertex  $i$  on sphere  $S_1$ ) to  $D_i^k$  (the disk tangent to vertex  $i$  on sphere  $S_k$ ).

The number of cylinders from  $S_1$  to  $S_k$  equals the total number of vertices  $n$ . This is equivalent to adding  $n - 1$  handles to  $S_1$ . Since we do this for each value of  $k$  from 2 to  $\theta$ , the resulting surface is a sphere with  $(\theta - 1)(n - 1)$  handles. We now erase the vertices on all spheres except  $S_1$  and extend incident edges on  $S_k$ ,  $k > 1$ , along the

incident cylinder to the corresponding vertex on  $S_1$ . (An equivalent way to view this pasting operation is that we are simply identifying all points of the boundaries of the two disks  $D_i^{1,k}$  and  $D_i^k$ , for  $i = 1, \dots, n$  and  $k = 2, \dots, \theta$ .)

This gives an embedding of  $G$  on a surface of genus  $(\theta - 1)(n - 1)$ , and therefore,  $\chi(G) \leq (\theta(G) - 1)(n - 1)$ .  $\square$

Since  $\theta \leq \lceil bt/2 \rceil$ , we obtain the following bound on  $\gamma$  in terms of  $bt$  and  $n$ .

**Corollary 7** *If  $G$  is a simple graph of order  $n$ , then*

$$\chi(G) \leq (\lceil bt(G)/2 \rceil - 1)(n - 1).$$

Since  $\theta \leq a_1$ , Theorem 7 also implies that  $\chi(G) \leq (a_1(G) - 1)(n - 1)$ . Using Euler's formula, we can cut this bound in half:

**Theorem 8** *If  $G$  is a simple graph with  $n$  vertices, then*

$$\chi(G) \leq (a_1(G) - 1)(n - 1) / 2.$$

**Proof** Suppose  $G$  can be written as a union of  $a_1$  forests, each on  $n$  vertices. Then  $e \leq a_1(n - 1)$ . If  $G$  has a 2-cell embedding on a surface of genus  $\gamma$ , then Euler's formula implies

$$\begin{aligned} 2 - 2\gamma &= n - e + f \geq n - a_1(n - 1) + f, \text{ and so} \\ 2\gamma &\leq (a_1 - 1)(n - 1) - f + 1, \end{aligned}$$

where  $f$  is the number of faces. Since  $f \geq 1$ , the result follows.  $\square$

Note that this result, together with the fact that  $a_1 \leq bt + 1$ , gives another bound on  $\gamma$  in terms of  $bt$  and  $n$ , almost identical to, but slightly weaker than, Corollary 7.

### 5. Conclusions

There are several places where the results in this work could be sharpened or improved. First, the examples known to the authors suggest that the constants can be lowered in Theorems 3, 7, and 8, and in Corollary 7. Second, it would be of interest to obtain a characterization of the graphs that satisfy  $\theta = O(\gamma/n)$ . Third, Theorem 5 would be improved by replacing the class of "short edge" graphs with another, more natural class of graphs of genus  $\gamma$ , for which  $a_1$  and  $\theta$  are bounded by a constant.

Finally, our work contains no upper bounds on book thickness. Yannakakis has proved a best possible result that  $bt \leq 4$  for  $G$  planar (cf. [Y]). Heath and Istrail use that result in [HI] to show that  $bt = O(\gamma)$ ; most recently, Malitz [M] has obtained a result of best possible order by showing that  $bt = O(\sqrt{\gamma})$ .

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