

Hyperelliptic Curve Arithmetic

Renate Scheidler



UNIVERSITY OF
CALGARY

12th International Conference on Finite Fields and Their
Applications
July 16, 2015

- **Number Theory**

- ▶ Invariant computation (class group/Jacobian, regulator, ...)
- ▶ Function field construction
- ▶ Function field tabulation

- **Geometry**

- ▶ Algebraic curves

- **Cryptography**

- ▶ Discrete log based crypto
- ▶ Pairing based crypto

- **Coding Theory?**

Groups that are used for **discrete log based crypto** should satisfy the following properties:

For practicality:

- Compact group elements
- Fast group operation

For security:

- Large order
- Cyclic or almost cyclic (plus some other restrictions on the order)
- Intractable discrete logarithm problem (DLP)

Proposed Groups:

- $G = \mathbb{F}_p^*$ (Diffie-Hellman 1976)
- Elliptic curves (Koblitz 1985, Miller 1985)
- Hyperelliptic curves (Koblitz 1989)

Fastest *generic* DLP algorithms: $O(\sqrt{|G|})$ group operations

- Best known for elliptic (i.e. genus 1) and genus 2 hyperelliptic curves
- Faster algorithms known for finite fields and higher genus curves

For curves of genus g over a finite field \mathbb{F}_q : $|G| \sim q^g$ as $q \rightarrow \infty$.

If we want 80 bits of security (i.e. $\sqrt{q^g} \approx 2^{80}$):

- $g = 1$: $q \approx 2^{160}$
- $g = 2$: $q \approx 2^{80}$ (slower group arithmetic but faster field arithmetic)

Let K be a field (in crypto, $K = \mathbb{F}_q$ with q prime or $q = 2^n$)

Weierstraß equation over K :

$$E : y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6 \quad (*)$$

with $a_1, a_2, a_3, a_4, a_6 \in K$

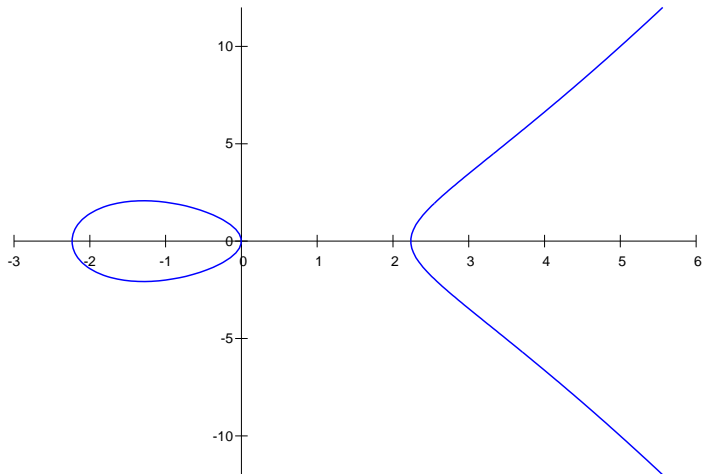
Elliptic curve: Weierstraß equation & **non-singularity** condition:
there are no simultaneous solutions to $(*)$ and

$$\begin{aligned} 2y + a_1x + a_3 &= 0 \\ a_1y &= 3x^2 + 2a_2x + a_4 \end{aligned}$$

Non-singularity $\iff \Delta \neq 0$ where Δ is the **discriminant** of E

An Example

$$E : y^2 = x^3 - 5x \text{ over } \mathbb{Q}$$



For $\text{char}(K) \neq 2, 3$, the variable transformations

$$y \rightarrow y - (a_1x + a_3)/2, \text{ then } x \rightarrow x - (a_1^2 + 4a_2)/12$$

yield an elliptic curve in **short** Weierstraß form:

$$E : y^2 = x^3 + Ax + B \quad (A, B \in K)$$

$$\text{Discriminant } \Delta = 4A^3 + 27B^2 \neq 0 \quad (\text{cubic in } x \text{ has distinct roots})$$

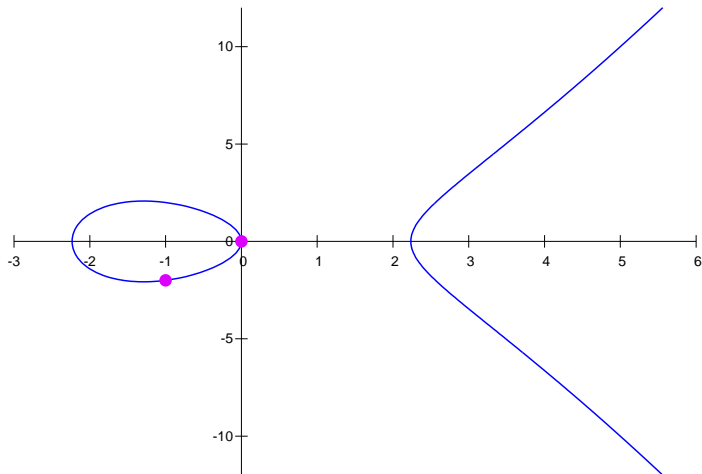
For any field L with $K \subseteq L \subseteq \overline{K}$:

$$E(L) = \{(x_0, y_0) \in L \times L \mid y_0^2 = x_0^3 + Ax_0 + B\} \cup \{\infty\}$$

set of **L -rational points** on E .

An Example

$$P_1 = (-1, 2), P_2 = (0, 0) \in E(\mathbb{Q})$$



In E , replace x by x/z , y by y/z , then multiply by z^3 :

$$E_{\text{proj}} : y^2z = x^3 + Axz^2 + Bz^3 .$$

Points on E_{proj} :

$[x : y : z] \neq [0 : 0 : 0]$, normalized so the last non-zero entry is 1.

Affine Points

Projective Points

$$(x, y) \leftrightarrow [x : y : 1]$$

$$\infty \leftrightarrow [0 : 1 : 0]$$

Goal: Make $E(K)$ into an additive (Abelian) group:

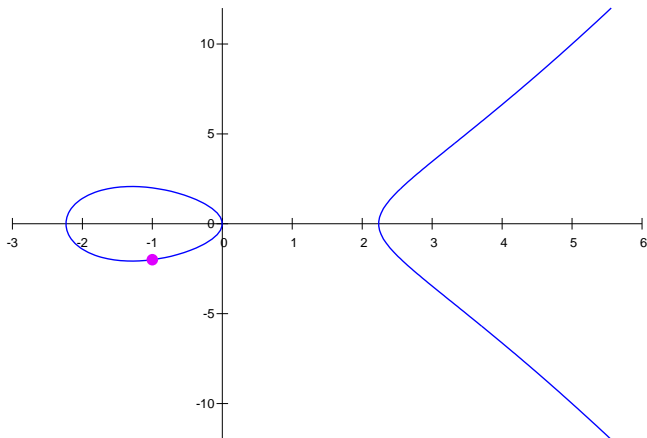
- The identity is the point at infinity.
- The inverse of a point $P = (x_0, y_0)$ is its **opposite** $\bar{P} = (x_0, -y_0)$ [†]
[†]true for odd characteristic only; in general, the opposite of a point $P = (x_0, y_0)$ is $\bar{P} = (x_0, -y_0 - a_1x_0 - a_3)$.

By **Bezout's Theorem**, any line intersects E in three points.

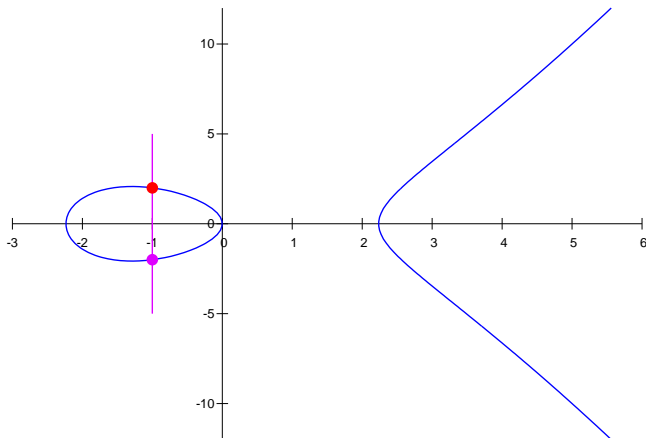
- Need to count multiplicities;
- If one of the points is ∞ , the line is “vertical”[†]
[†]true for odd characteristic only; in general, the line goes through P and \bar{P} .

Motto: “Any three collinear points on E sum to zero (i.e. ∞).”

Also known as **Chord & Tangent Addition Law**.

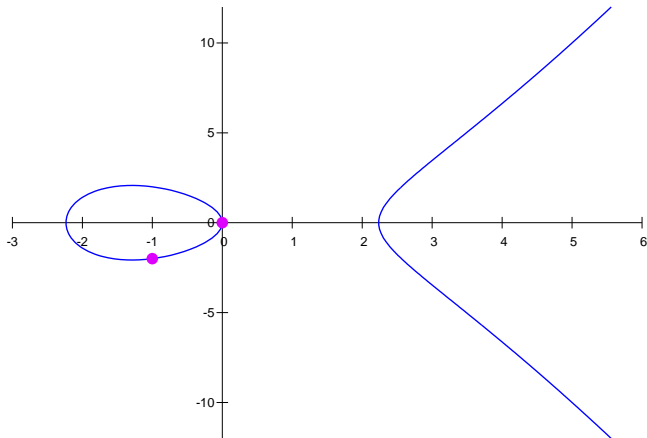


$$-(\bullet) = ?$$



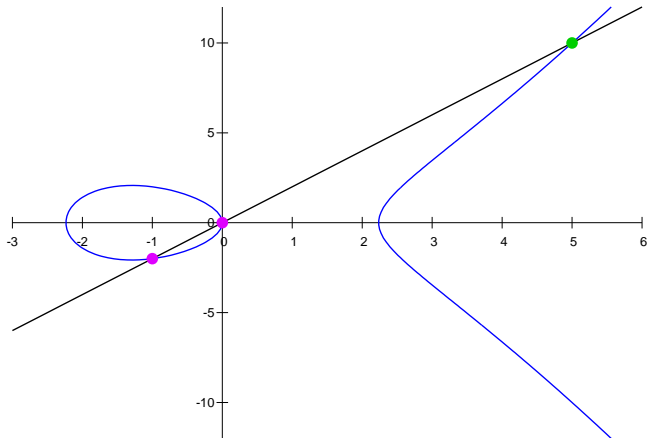
$$\bullet + \bullet + \infty = 0 \quad \Rightarrow \quad -(\bullet) = \bullet$$

Addition on Elliptic Curves



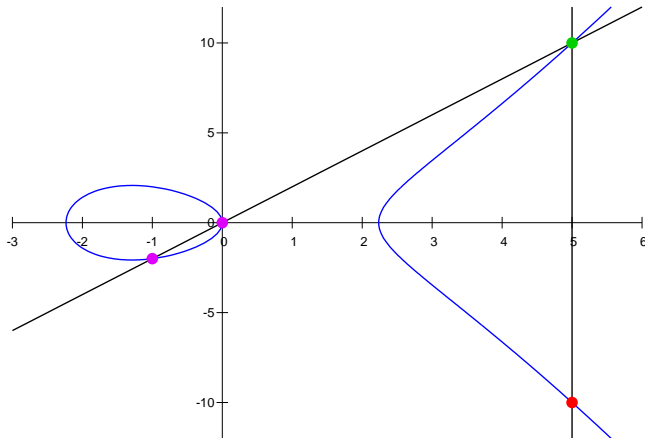
$$\bullet + \bullet = ?$$

Addition on Elliptic Curves



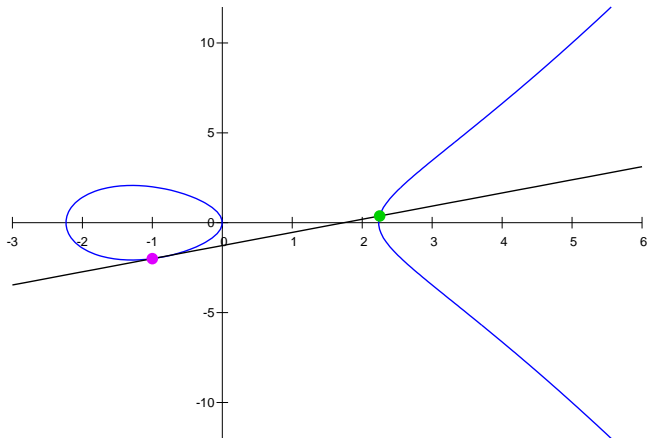
$$\bullet + \bullet + \bullet = 0$$

Addition on Elliptic Curves



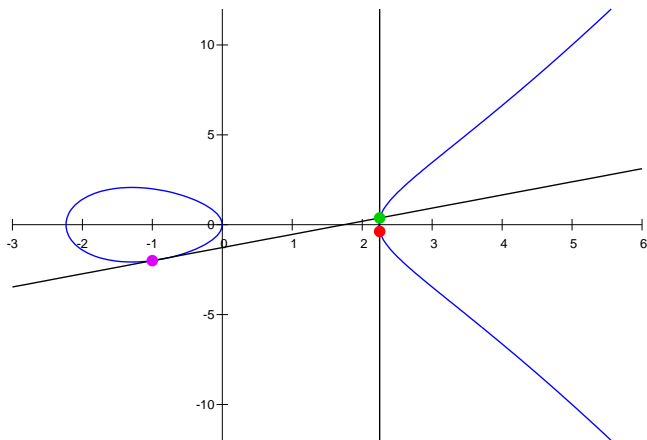
$$\bullet + \bullet + \bullet = 0 \quad \Rightarrow \quad \bullet + \bullet = \bullet$$

Doubling on Elliptic Curves



$$2 \times \bullet = ?$$

Doubling on Elliptic Curves



$$2 \times \text{pink} + \text{green} = 0 \quad \Rightarrow \quad 2 \times \text{pink} = \text{red}$$

Let

$$P_1 = (x_1, y_1), P_2 = (x_2, y_2) \quad (P_1 \neq \infty, P_2 \neq \infty, P_1 + P_2 \neq \infty).$$

Then

$$-P_1 = (-x_1, y_1)$$

$$P_1 + P_2 = (\lambda^2 - x_1 - x_2, -\lambda^3 + \lambda(x_1 + x_2) - \mu)$$

where

$$\lambda = \begin{cases} \frac{y_2 - y_1}{x_2 - x_1} & \text{if } P_1 \neq P_2 \\ \frac{3x_1^2 + A}{2y_1} & \text{if } P_1 = P_2 \end{cases}$$

$$\mu = \begin{cases} \frac{y_1x_2 - y_2x_1}{x_2 - x_1} & \text{if } P_1 \neq P_2 \\ \frac{-x_1^3 + Ax_1 + 2B}{2y_1} & \text{if } P_1 = P_2 \end{cases}$$

Recall Weierstraß equation:

$$E : y^2 + \underbrace{(a_1x + a_3)}_{h(x)}y = \underbrace{x^3 + a_2x^2 + a_4x + a_6}_{f(x)}$$

$$\deg(f) = 3 = 2 \cdot 1 + 1 \text{ odd}$$

$$\deg(h) = 1 \text{ for } \text{char}(K) = 2; \quad h = 0 \text{ for } \text{char}(K) \neq 2$$

$$\text{Generalization: } \deg(f) = 2g + 1, \quad \deg(h) \leq g$$

g is the **genus** of the curve

$g = 1$: elliptic curves

$g = 2$: $\deg(f) = 5, \deg(h) \leq 2$ (always hyperelliptic).

Hyperelliptic curve of genus g over K :

$$H : y^2 + h(x)y = f(x)$$

- $h(x), f(x) \in K[x]$
- $f(x)$ monic and $\deg(f) = 2g + 1$ is odd
- $\deg(h) \leq g$ if $\text{char}(K) = 2$; $h(x) = 0$ if $\text{char}(K) \neq 2$
- non-singularity

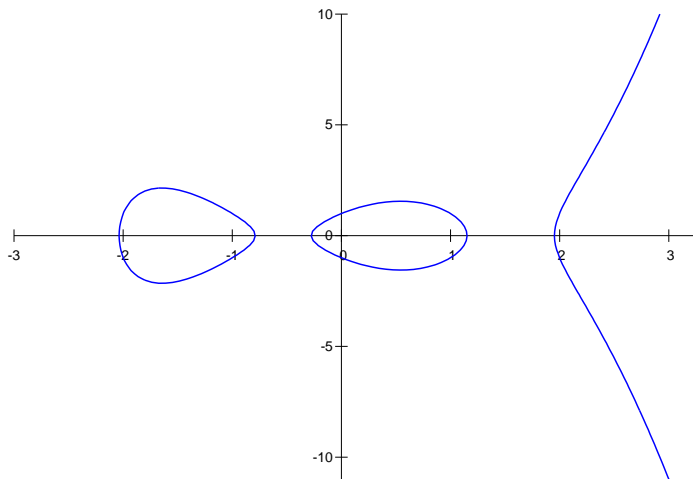
$\text{char}(K) \neq 2$: $y^2 = f(x)$, $f(x)$ monic, of odd degree, square-free

Set of L -rational points on H ($K \subseteq L \subseteq \bar{K}$):

$$H(L) = \{(x_0, y_0) \in L \times L \mid y_0^2 + h(x_0)y_0 = f(x_0)\} \cup \{\infty\}$$

An Example

$H : y^2 = x^5 - 5x^3 + 4x - 1$ over \mathbb{Q} , genus $g = 2$



- Group of **divisors** on H :

$$\text{Div}_H(\bar{K}) = \langle H(\bar{K}) \rangle = \left\{ \sum_{\text{finite}} m_P P \mid m_P \in \mathbb{Z}, P \in H(\bar{K}) \right\}$$

- Subgroup of $\text{Div}_H(\bar{K})$ of **degree zero divisors** on H :

$$\text{Div}_H^0(\bar{K}) = \langle [P] \mid P \in H(\bar{K}) \rangle = \left\{ \sum_{\text{finite}} m_P [P] \mid m_P \in \mathbb{Z}, P \in H(\bar{K}) \right\}$$

where $[P] = P - \infty$

- Subgroup of $\text{Div}_H^0(\bar{K})$ of **principal divisors** on H :

$$\text{Prin}_H(\bar{K}) = \left\{ \sum_{\text{finite}} v_P(\alpha) [P] \mid \alpha \in K(x, y), P \in H(\bar{K}) \right\}$$

Jacobian of H : $\text{Jac}_H(\bar{K}) = \text{Div}_H^0(\bar{K}) / \text{Prin}_H(\bar{K})$

Motto: “Any complete collection of points on a *function* sums to zero.”

$$H(\bar{K}) \hookrightarrow \text{Jac}_H(\bar{K}) \quad \text{via } P \mapsto [P]$$

For elliptic curves: $E(\bar{K}) \cong \text{Jac}_E(\bar{K}) \quad (\Rightarrow E(\bar{K}) \text{ is a group})$

Identity: $[\infty] = \infty - \infty$

Inverses: The points

$$P = (x_0, y_0) \quad \text{and} \quad \bar{P} = (x_0, -y_0 - h(x_0))$$

on H both lie on the function $x = x_0$, so

$$-[P] = [\bar{P}]$$

Every class in $\text{Jac}_H(\overline{K})$ contains a divisor $\sum_{\text{finite}} m_P [P]$ such that

- all $m_P > 0$ (replace $-[P]$ by $[\overline{P}]$)
- if $P = \overline{P}$, then $m_P = 1$ (as $2[P] = 0$)
- if $P \neq \overline{P}$, then only one of P, \overline{P} can appear in the sum (as $[P] + [\overline{P}] = 0$)

Such a divisor is **semi-reduced**. If $\sum m_P \leq g$, then it is **reduced**.

E.g. $g = 2$: reduced divisors are of the form $[P]$ or $[P] + [Q]$.

Theorem

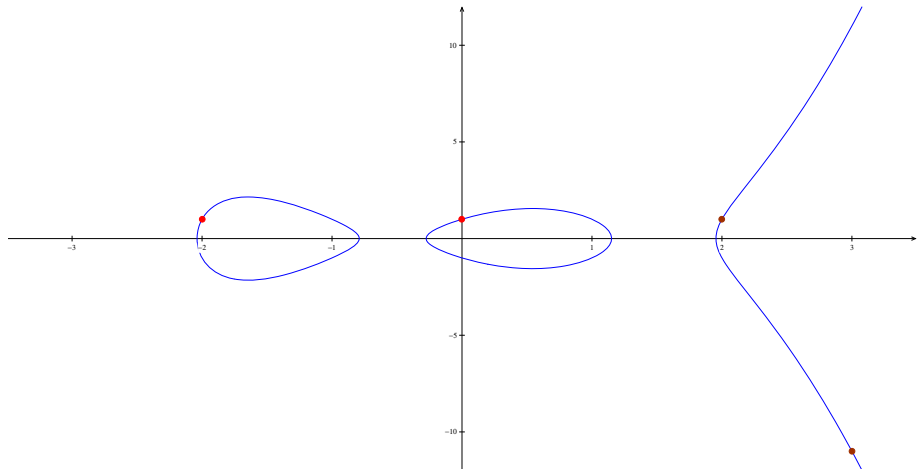
Every class in $\text{Jac}_H(\overline{K})$ contains a unique reduced divisor.

For reduced D_1, D_2 , the reduced divisor in the class $[D_1 + D_2]$ is denoted $D_1 \oplus D_2$.

An Example of Reduced Divisors

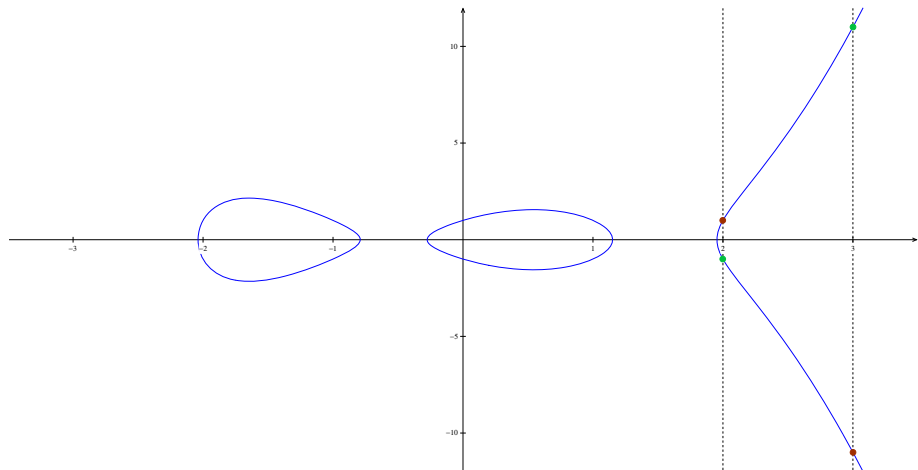
$$D_1 = (-2, 1) + (0, 1)$$

$$D_2 = (2, 1) + (3, -11)$$



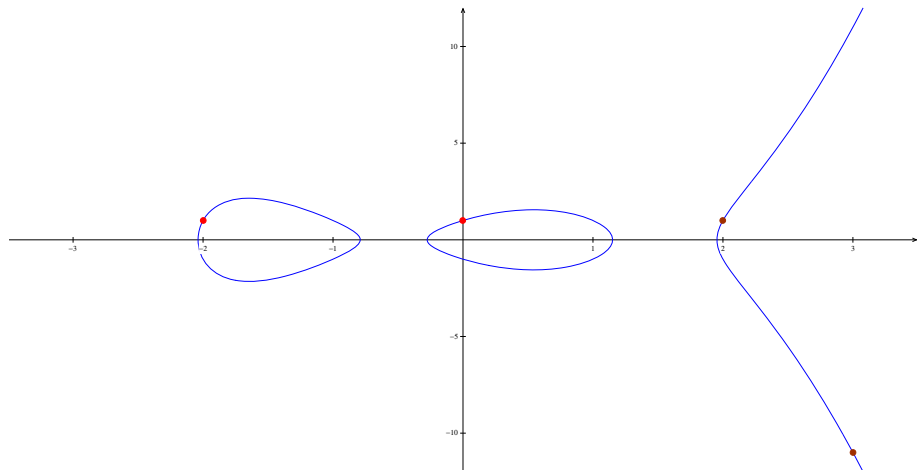
Inverses on Hyperelliptic Curves

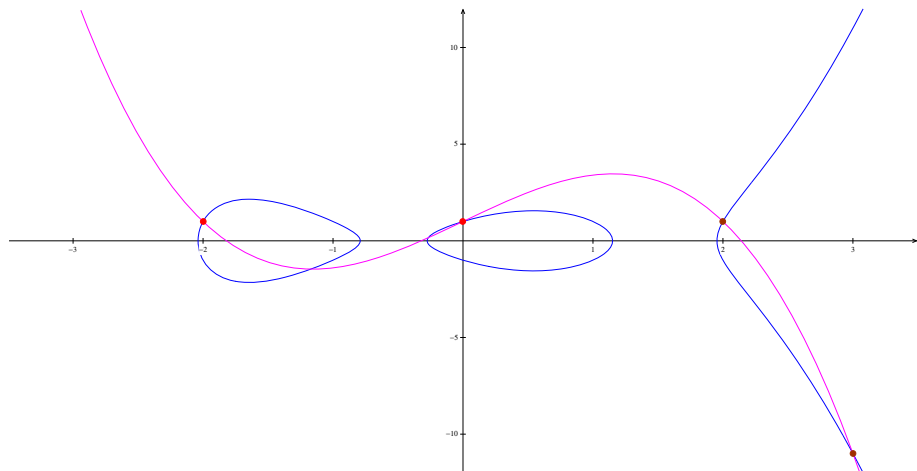
The inverse of $D = P_1 + P_2 + \cdots + P_r$ is $-D = \bar{P}_1 + \bar{P}_2 + \cdots + \bar{P}_r$

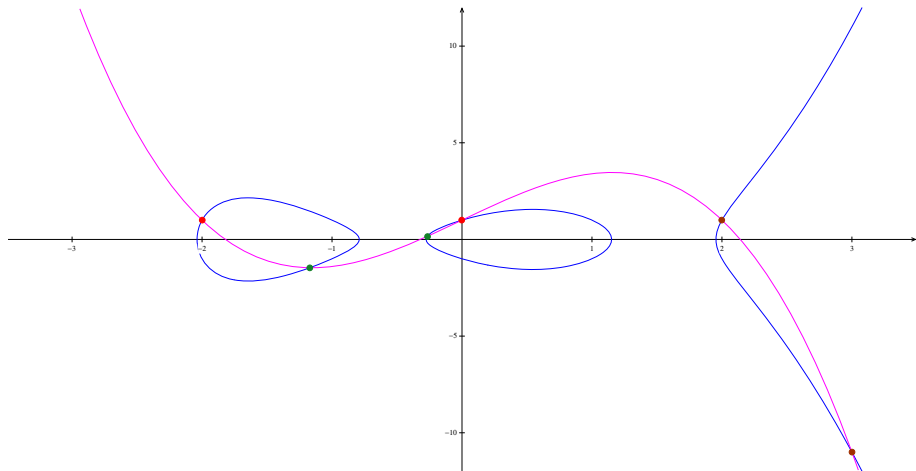


$$-(\bullet + \bullet) = (\bullet + \bullet)$$

Addition on Genus 2 Curves

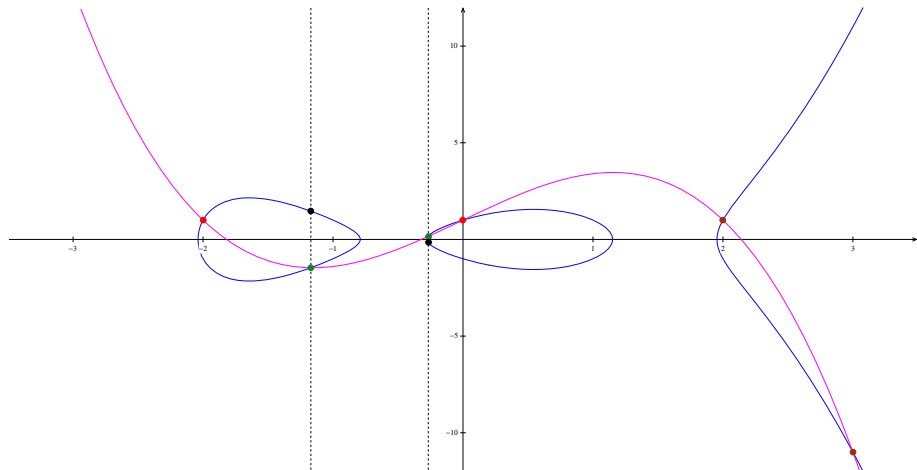






$$(\bullet + \bullet) + (\bullet + \bullet) + (\bullet + \bullet) = 0$$

Addition on Genus 2 Curves



$$(\bullet + \bullet) + (\bullet + \bullet) + (\bullet + \bullet) = 0 \quad \Rightarrow \quad (\bullet + \bullet) \oplus (\bullet + \bullet) = (\bullet + \bullet)$$

Motto: “Any complete collection of points on a function sums to zero.”

To add and reduce two divisors $P_1 + P_2$ and $Q_1 + Q_2$ in genus 2:

- The four points P_1, P_2, Q_1, Q_2 lie on a unique function $y = v(x)$ with $\deg(v) = 3$.
- This function intersects H in two more points R_1 and R_2 :
 - ▶ The x -coordinates of R_1 and R_2 can be obtained by finding the remaining two roots of $v(x)^2 + h(x)v(x) = f(x)$.
 - ▶ The y -coordinates of R_1 and R_2 can be obtained by substituting the x -coordinates into $y = v(x)$.
- Since $(P_1 + P_2) + (Q_1 + Q_2) + (R_1 + R_2) = 0$, we have

$$(P_1 + P_2) \oplus (Q_1 + Q_2) = \overline{R_1} + \overline{R_2} .$$

Addition in Genus 2 – Example

Consider $H : y^2 = f(x)$ with $f(x) = x^5 - 5x^3 + 4x + 1$ over \mathbb{Q} .

To add & reduce $(-2, 1) + (0, 1)$ and $(2, 1) + (3, -11)$, proceed as follows:

- The unique degree 3 function through $(-2, 1)$, $(0, 1)$, $(2, 1)$ and $(3, -11)$ is $y = v(x)$ with $v(x) = -(4/5)x^3 + (16/5)x + 1$.
- The equation $v(x)^2 = f(x)$ becomes

$$(x - (-2))(x - 0)(x - 2)(x - 3)(16x^2 + 23x + 5) = 0 .$$

- The roots of $16x^2 + 23x + 5$ are $\frac{-23 \pm \sqrt{209}}{32}$.
- The corresponding y -coordinates are $\frac{-1333 \pm 115\sqrt{209}}{2048}$. So

$$\begin{aligned} &(-2, 1) + (0, 1) \oplus (2, 1) + (3, -11) = \\ &\left(\frac{-23 + \sqrt{209}}{32}, \frac{1333 - 115\sqrt{209}}{2048} \right) + \left(\frac{-23 - \sqrt{209}}{32}, \frac{1333 + 115\sqrt{209}}{2048} \right). \end{aligned}$$

Let D_1, D_2 be reduced divisors on $H : y^2 + h(x)y = f(x)$.

First form the **semi-reduced sum** of D_1 and D_2 , obtaining $D = \sum_{i=1}^r [P_i]$

Now iterate over D as follows, until $r \leq g$:

- The r points P_i all lie on a curve $y = v(x)$ with $\deg(v) = r - 1$.
- $w(x) = v^2 - hv - f$ is a polynomial of degree $\max\{2r - 2, 2g + 1\}$.
 r of the roots of $w(x)$ are the x -coordinates of the P_i .
- If $r \geq g + 2$, then $\deg(w) = 2r - 2$, yielding $r - 2$ further roots.
If $r = g + 1$, then $\deg(w) = 2g + 1$, yielding g further roots.
- Substitute these new roots into $y = v(x)$ to obtain $\max\{r - 2, g\}$ new points on H . Replace D by the new divisor thus obtained.

Since $r \leq 2g$ at the start, $D_1 \oplus D_2$ is obtained after at most $\lceil g/2 \rceil$ steps.

Let $D = \sum_{i=1}^r m_i [P_i]$ be a semi-reduced divisor, $P_i = (x_i, y_i)$

The **Mumford representation** of D is a pair of polynomials $(u(x), v(x))$ that uniquely determines D :

$u(x)$ captures all the x -coordinates with multiplicities;

$y = v(x)$ is the interpolation function through all the P_i (as before).

Formally:

$$u(x) = \prod_{i=1}^r (x - x_i)^{m_i}$$

$$\left(\frac{d}{dx}\right)^j [v(x)^2 + v(x)h(x) - f(x)]_{x=x_i} = 0 \quad (0 \leq j \leq m_i - 1)$$

Properties:

- $u(x_i) = 0$ and $v(x_i) = y_i$ with multiplicity m_i for $1 \leq i \leq r$;
- $u(x)$ is monic and divides $v(x)^2 + h(x)v(x) - f(x)$
- D uniquely determines $u(x)$ and $v(x) \bmod u(x)$;
- Any pair of polynomials $u(x), v(x) \in \overline{K}[x]$ with $u(x)$ monic and dividing $v(x)^2 + h(x)v(x) - f(x)$ determines a semi-reduced divisor.

Examples:

- If $D = [(x_0, y_0)]$ is a point, then $u(x) = x - x_0$ and $v(x) = y_0$.
- If $D = [(x_1, y_1)] \oplus [(x_2, y_2)]$, then
$$u(x) = (x - x_1)(x - x_2),$$
$$y = v(x) \text{ is the line through } (x_1, y_1) \text{ and } (x_2, y_2).$$

Let $D_1 = (u_1, v_1)$, $D_2 = (u_2, v_2)$.

Simplest case: for any $[P]$ occurring in D_1 , $[\bar{P}]$ doesn't occur in D_2 and vice versa. Then $D_1 + D_2 = (u, v)$ is semi-reduced and

$$u = u_1 u_2, \quad v = \begin{cases} v_1 & (\text{mod } u_1), \\ v_2 & (\text{mod } u_2). \end{cases}$$

In general: suppose $P = (x_0, y_0)$ occurs in D_1 and \bar{P} occurs in D_2 .

Then $u_1(x_0) = u_2(x_0) = 0$ and $v_1(x_0) = y_0 = -v_2(x_0) - h(x_0)$, so $x - x_0$ divides $u_1(x)$, $u_2(x)$, $v_1(x) + v_2(x) + h(x)$.

$$d = \gcd(u_1, u_2, v_1 + v_2 + h) = s_1 u_1 + s_2 u_2 + s_3 (v_1 + v_2 + h).$$

$$u = u_1 u_2 / d^2.$$

$$v \equiv \frac{1}{d} (s_1 u_1 v_2 + s_2 u_2 v_1 + s_3 (v_1 v_2 + f)) \pmod{u}$$

(In the simplest case above, $d = 1$ and $s_3 = 0$)

Let $D = (u, v)$ be a semi-reduced divisor on $H : y^2 + h(x)y = f(x)$.

While $\deg(u) > g$ do

// Replace the x -coordinates of the points in D by those of the other intersection points of H with v :

$$u \leftarrow (f - vh - v^2)/u \ .$$

// Replace the new points by their opposites:

$$v \leftarrow (-v - h) \pmod{u} \ .$$

Consider again $H : y^2 = f(x)$ with $f(x) = x^5 - 5x^3 + 4x + 1$ over \mathbb{Q} .

Compute $D_1 \oplus D_2$ with $D_1 = (-2, 1) + (0, 1)$ and $D_2 = (2, 1) + (3, -11)$:

Mumford rep of D_1 : $u_1(x) = x^2 + 2x$, $v_1(x) = 1$.

Mumford rep of D_2 : $u_2(x) = x^2 - 5x + 6$, $v_2(x) = -12x + 25$.

$$u(x) = u_1(x)u_2(x) = x^4 - 3x^3 - 4x^2 + 12x ;$$

$$v(x) = -(4/5)x^3 + (16/5)x + 1 ;$$

$$u(x) \leftarrow (f(x) - v(x)^2)/u(x) = 16x^2 + 23x + 5 ;$$

$$v \leftarrow -v \pmod{u} = (16x - 23)/320 ;$$

Mumford rep of $D_1 \oplus D_2 = \left(\frac{-23 + \sqrt{209}}{32}, \frac{1333 - 115\sqrt{209}}{2048} \right) + \left(\frac{-23 - \sqrt{209}}{32}, \frac{1333 + 115\sqrt{209}}{2048} \right)$:

$$u(x) = 16x^2 + 23x + 5, \quad v(x) = (16x - 23)/320.$$

Divisors defined over K

Let $\phi \in \text{Gal}(\overline{K}/K)$ (for $K = \mathbb{F}_q$, think of Frobenius $\phi(\alpha) = \alpha^q$).

ϕ acts on points via their coordinates, and on divisors via their points.

A divisor D is **defined over** K if $\phi(D) = D$ for all $\phi \in \text{Gal}(\overline{K}/K)$.

Example: The divisor

$$D = \left(\frac{-23 + \sqrt{209}}{32}, \frac{1333 - 115\sqrt{209}}{2048} \right) + \left(\frac{-23 - \sqrt{209}}{32}, \frac{1333 + 115\sqrt{209}}{2048} \right)$$

is defined over \mathbb{Q} (invariant under automorphism $\sqrt{209} \mapsto -\sqrt{209}$).

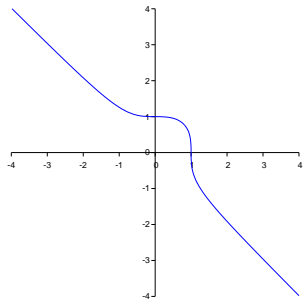
Theorem

$D = (u, v)$ is defined over K if and only if $u(x), v(x) \in K[x]$.

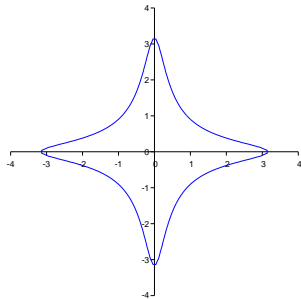
Corollary

If K is a finite field, then $\text{Jac}(H)$ is finite.

- **Hessians:** $x^3 + y^3 - 3dxy = 1$
- **Edwards models:** $x^2 + y^2 = c^2(1 + dx^2y^2)$ (q odd) and variations

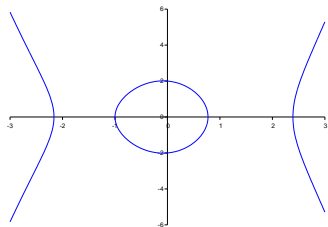


$$x^3 + y^3 = 1$$



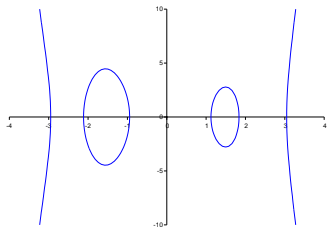
$$x^2 + y^2 = 10(1 - x^2y^2)$$

$y^2 + h(x)y = f(x)$, $\deg(f) = 2g + 2$, $\deg(h) = g + 1$ if $\text{char}(K) = 2$.



$$y^2 = x^4 - 6x^2 + x + 6$$

$(g = 1)$



$$y^2 = x^6 - 13x^4 + 44x^2 - 4x - 1$$

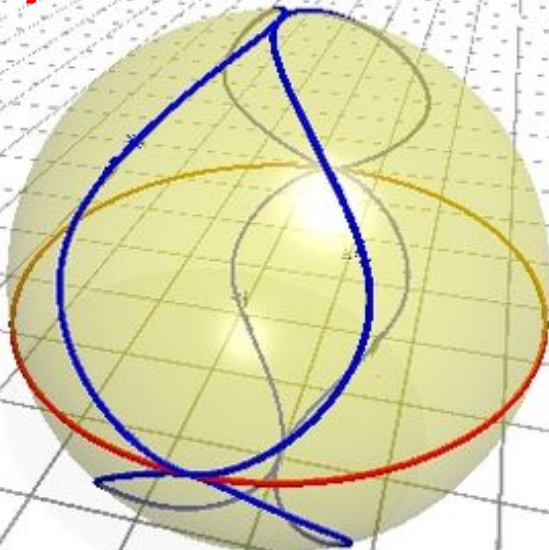
$(g = 2)$

- More general and plentiful than odd degree hyperelliptic curves:
 - ▶ can always transform an odd to even degree model over K , but the reverse direction may require an extension of K .
- Two points at infinity (∞ and $\overline{\infty}$).
- Divisor Representation: $\sum_{i=1}^r P_i - r\infty + n(\overline{\infty} - \infty)$, $r \leq g$.
 - ▶ No restrictions on n : many reduced divisors in each class ($\approx q^g$)
 - ▶ $n = 0$: **infrastructures** (misses a few divisor classes)
 - ▶ $n \approx g$: unique representatives (Paulus-Rück 1999)
 - ▶ $n \approx \lceil g/2 \rceil$: **balanced representation**, unique and much better for computation (Galbraith-Harrison-Mireles Morales 2008)
- The DLP's in all these settings are polynomially equivalent.

- Genus 1 and 2, q prime or $q = 2^n$: efficient and secure for DLP based crypto. Genus 3 might also be OK.
- **Explicit formulas** reduce the polynomial arithmetic to arithmetic in \mathbb{F}_q .
Odd degree: LOTS of literature on genus 2, a bit on genus 3 and 4;
Even degree: reasonably developed for genus 2, work on genus 3 in progress.
- Other coordinates (e.g. projective coordinates) can be more efficient. They avoid inversions in \mathbb{F}_q , at the expense of redundancy. Oftentimes *mixed* coordinates are best.
- For genus 1, use Edwards models — more efficient, *unified* formulas. No higher genus Edwards analogue is known.
- For genus 2 and odd degree, Gaudry's *Kummer surface* arithmetic is fastest, but doesn't work for all curves.
- Work on arbitrary genus is ongoing.

Thank you!

Questions?



$$y^2 = x^6 + x^2 + x$$